

# **Boolean Synthesis via Decomposition**

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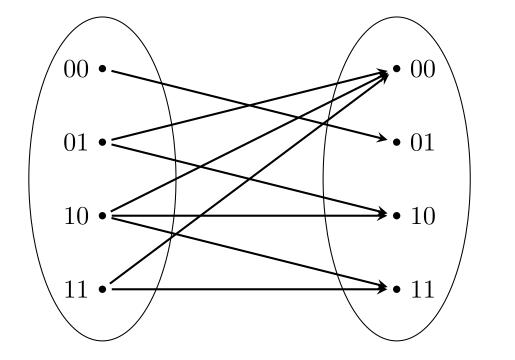
## Boolean Synthesis [1]

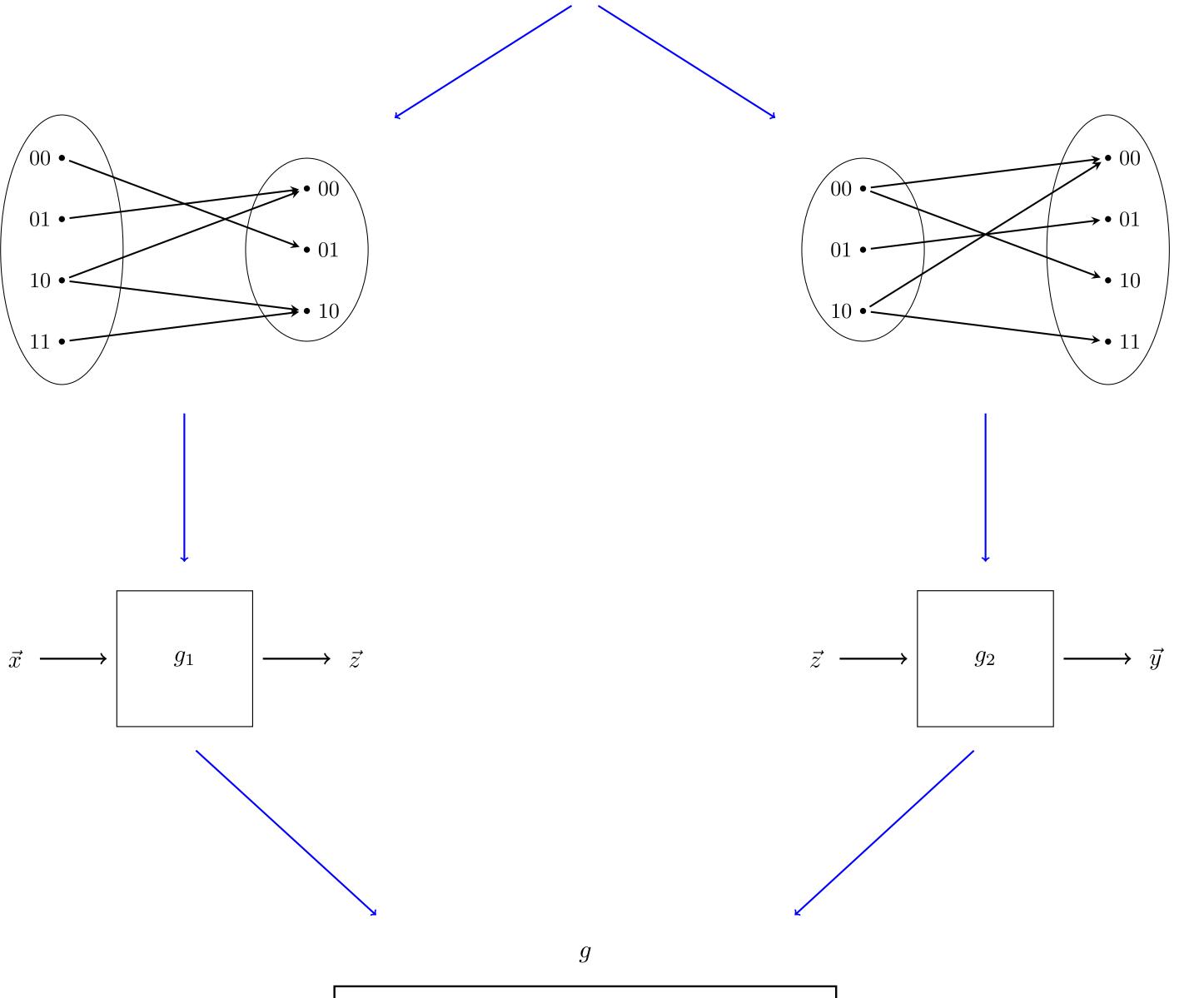
**Given:** Boolean formula  $F(\vec{x}, \vec{y})$  representing a relation over input variables  $\vec{x} = \{x_1, \ldots, x_m\}$  and output variables  $\vec{y} = \{y_1, \ldots, y_n\}$ 

**Obtain:** Boolean function  $g: \{0,1\}^m \to \{0,1\}^n$  such that, for all  $\vec{x}$ ,

#### **Towards Sequential Decomposition**

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### $F(\vec{x}, g(\vec{x})) \Leftrightarrow \exists \vec{y}. F(\vec{x}, \vec{y})$

• F is called the *specification*.

• g is called the *implementation*.

**Example:** The two's complement of a two-bit integer  $x_1x_0$  is a two-bit integer  $y_1y_0$  such that  $x_1x_0 + y_1y_0 = 0$ . We can synthesize a function that computes the two's complement as follows:

$$F(x_0, x_1, y_0, y_1) = \neg (x_0 \oplus y_0) \land \neg (x_1 \oplus y_1 \oplus (x_0 \land y_0))$$
 $\Downarrow$ 
 $g(x_0, x_1) = \begin{cases} y_0 := x_0 \\ y_1 := x_1 \oplus x_0 \end{cases}$ 

Despite extensive research on the subject, Boolean synthesis remains a challenging NP-hard problem.

A standard strategy for handling hard problems is decomposing them into smaller problems. Our goal is to apply this concept to Boolean synthesis.

# **Decomposition using Factored Formulas**

One way to decompose Boolean synthesis is to use factored formulas [2, 3]:

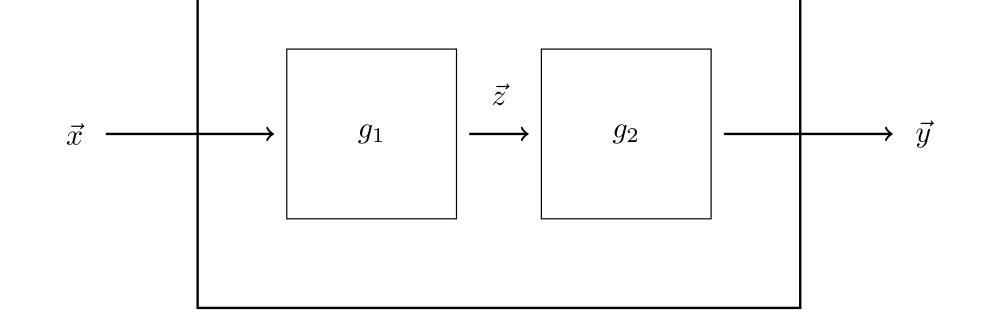
 $F(\vec{x}, y_1, y_2, y_3, y_4) = F_1(\vec{x}, y_2, y_4) \wedge F_2(\vec{x}, y_1, y_2, y_3) \wedge F_3(\vec{x}, y_3)$ 

#### **Pros:**

- Easy to perform decomposition.
- Specifications are often already given as a conjunction of constraints.
- Each factor uses only a subset of the variables.

#### Cons:

- Dependences between factors.
- Highly non-trivial to combine implementations of  $F_1, \ldots, F_k$  into an implementation of F [2].



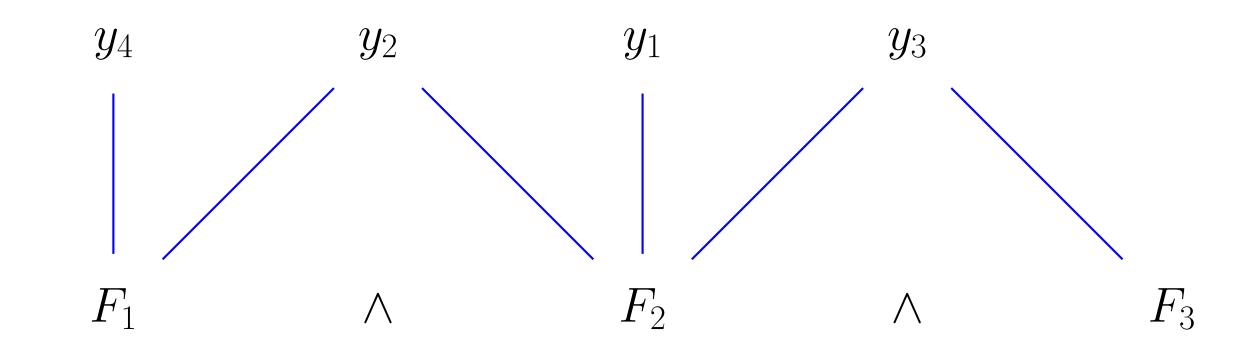
# Sequential Decomposition

**Given:** Boolean formula  $F(\vec{x}, \vec{y})$  representing a relation over input variables  $\vec{x} = \{x_1, \dots, x_m\}$  and output variables  $\vec{y} = \{y_1, \dots, y_n\}$ 

**Obtain:** Formulas  $F_1(\vec{x}, \vec{z})$  and  $F_2(\vec{z}, \vec{y})$  for intermediate variables  $\vec{z} = \{z_1, \ldots, z_k\}$  such that, if  $g_1$  implements  $F_1$  and  $g_2$  implements  $F_2$ , then  $g_2 \circ g_1$  implements F.



This form of decomposition has been shown to significantly improve synthesis algorithms [2, 3]. However, dealing with the dependences between factors prevents us from taking full advantage of the decomposition [3]:



[1] Dror Fried, Lucas M. Tabajara, and Moshe Y. Vardi. BDD-Based Boolean Functional Synthesis.

In Computer Aided Verification - 28th International Conference, CAV 2016, Toronto, ON, Canada, July 17-23, 2016, Proceedings, Part II. Springer, 2016.

[2] Ajith K. John, Shetal Shah, Supratik Chakraborty, Ashutosh Trivedi, and S. Akshay.
 Skolem Functions for Factored Formulas.
 In Formal Methods in Computer-Aided Design, FMCAD 2015, Austin, Texas,

USA, September 27-30, 2015., pages 73–80, 2015.

[3] Lucas M. Tabajara and Moshe Y. Vardi. Factored Boolean Functional Synthesis.

In Formal Methods in Computer-Aided Design, FMCAD 2017, Vienna, Austria, October 2-6, 2017., 2017.